



zIIP stealing GCP MSUs time for Capacity Management



Tuesday 7th November, 2017 (10:45 – 11:45) Session LI in Woodcote

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- Elimination of zAAP good or bad?
- Why should we care ? –
- Who are zIIP exploiters? –
- · How to get more zIIP capacity -
- How to control overflow to GCP -
- How to reduce zIIP usage -









Elimination of zAAP

Pros

- CP utilization Improved
- MP effect Reduced
- Queueing time Reduced w/ more eligible CPs
- MLC save w/ less hardware
 - Small zIIP & zAAP workload could justify an engine

Cons

- Application JAVA code contending with system background tasks
- Eliminated separate
 overflow controls
 - Workloads impacted by delays <u>were not</u> controllable



zAAP Demise

- zAAP on zIIP
 - z/OS V1.11
 - ZAAPZIIP=YES
 - z/OS V1.12
 - APAR to allow zAAP to run on zIIP when zAAP present

zEnterprise 196 & 114

zIIP to CP ratio <= 1:1

zEC12 and zBC12

- zIIP to CP ratio <= 2:1
- Last CPCs w/ zAAPs



Multiprocessor (MP) effect



- Original underutilized
 - 8 zIIP and
 - -8 zAAP

 If fewer CPs then lower MP effect and better utilization / CP



Why / how overflow



Defaults allow specialty engine to overflow to GCP

- Spare specialty of other type wasted
- Controls were LPAR wide
- zAAP had more controls
- Help signaled if <u>all busy</u> of 1 specialty type



zIIP % Busy and Crossover



- 1 Specialty engine
 @ 35%
- 2 Specialty CP @ 50-60%
- Only 1 type = less likely all busy to ask for Help





Why Should I Care?



Overloaded zIIPs cost GCP MSUs

- zIIPs overloaded w/o zAAP work
- Large MSUs overflowing to GCPs
 - Slowing down workloads
 - Costing hardware and software \$\$\$
- More zIIP exploiters on horizon



IMPACTING SLAs & Costs



Overflow MSU to GCPs = MLC \$

<u>MLC Impact</u> : \$250K – \$750K

per year ...per engine wasted



File Edit Graphics Database Tools Options Window Help





- 77 MSUs @ \$400 / MSU = \$372K / yr
- Smaller sites pay > \$/ MSU larger sites waste more **MSUs**





Who are zllP Exploiters?



Who is using my zIIP

Traditional zIIP Users

- DB2 v8+ & IMS v12+ (more to come)
- Batch (SyncSort)
- XML Sys Service, ComSrvr IPSec
- Global Mirror DFSMS Data Mover
- Performance monitors

zAAP on zIIP Users

- Batch JAVA (JZOS)
- Transactional JAVA
 - WAS, CICS Liberty, IMS JMP, Spark, etc.



Rock 'Em Sock 'Em Robots ©Mattel



DB2 zIIP Eligible

- **DB2 buffer pool processing** up to 100% for
 - prefetch, Deferred write, Page set castout (CF Read GBP),
- DB2 Utilities up to 100% of
 - LOAD, REORG, and REBUILD INDEX
- Log processing read, and write
- DRDA over IP workloads SQL request up to 60%
- Parallel query child processes up to 80%
- XML Schema processing up to 100% of
 - schema validation and non-validation parsing



IMS zIIP Eligible

IMS V12 – CQS access ShrMsgQs IMS V13

- **IMS Connect** thread requests via TCP/IP for;
 - SOAP messages
 - DRDA®
 - MSC
 - ISC

ODBM address space

- DRDA® threads requests via TCP/IP
- Threads requests via CSLDMI API.

zAAP on zIIP

- Java Message Processing (JMP) regions
- Java Batch Processing (JBP) regions
- Message Processing Programs (MPP) w/ COBOL to JAVA



Harder to tell what is using zIIP?



- 55% to DB2
 - Prefetch
 - Utilities
- 35% to Java
 - Application
 - Garbage
 Collection
- 2% to Monitors
- 3% to Overhead
- 5% spare?





How to get more zIIP capacity?

Preventing Overflow – Increase Capacity

- Buy more zIIP or can't due to
 - GCP Ratio restrictions
 - Hardware \$s
- Other Options
 - Sub-capacity CPCs zIIP MSUs to GCP ratio better
 - Simultaneous Multithreading (SMT) on z13+ family





Subcap – zIIP Relative Speed

z13	Model	MSUs	2 zIIP	Eff Ratio
2964	401	31	420	13.55
2964	501	94	420	4.47
2964	601	134	420	3.13
2964	701	210	420	2

- zIIP, zAAP run at full 701 speed on all models
- 2 zIIPs / GCP can be 13.5x more MSUs instead of 2x on 7xx models



Cost of Subcapacity GCP crossover

- Is cross-over occurring in 4HRA CPC peak?
- Higher for Sub-capacity model w/ GCPs slower than zIIPs
 - -10% overflow of one zIIP = 21 MSUs or 68% of 1w 401 GCP
 - 87% of 1 CP on 2964-416
 - Normalized zIIP Seconds in Monitors & SMF Recs
 - Address space(30), Processor(70), SrvCls(72)

z13	Model	MSUs	1 zIIP	Eff Ratio
2964	401	31	210	6.77

z13 Simultaneous MultiThreading (SMT)



to pass

GUIDE



SMT Effectiveness



Capacity Factor

Any CF # > 1
 means SMT is
 delivering
 additional zIIP
 capacity



zIIP Overflow



 Overflow still occurring, but would have been more MSUs



Overflow w/o > 65+%



- 44% Busy
 - 4 of 9 zIIP busy
- Queueing theory
 - predicts elapsed time elongation based on # of servers / zIIPs
 - > 35% for 1 is begin of CP queueing
 - Does not predict probability of all busy for 1 microsec



SMT2 – not at Max Capacity



 Just like SMT1 mode most work does not wait for zIIP if all zIIPs busy at the moment in time



SMT – Other benefits

- zIIP initiated CF SYNC calls w/ SMT do not spin both threads
 - Hiperlinks
 - DB2 Utility page locks





How to control overflow to GCP?



Specialty CP on GCP Controls

IEAOPTxx Options

Crossover	Honor Priority	General Purpose Processors Behavior
NO	VEC	zIIP / zAAP work on GCPs only
		when help needed
VEC		zAAP work on GCPs only when
TES	NU	no non-zAAP work is ready (1)
		zAAP work on GCPs when help
YES	<u>YES</u>	needed and no non-zAAP work
		ready
<u>NO</u>	NO	No zIIP / zAAP work on GCPs (1)

1) HonorPrio=NO; Specialty engine eligible only runs on GCPs to resolve resource conflicts w/ GCP work

Help needed is signaled when all specialty CPs of a type are busy. Less likely if all 1 type.

zAAP

- IFACROSSOVER
- IFAHONORPRIORITY - Parm z/OS 1.6 (2005)

zIIP

- $zIIP Crossover = \underline{NO}$
 - GCPs help if zIIP asks
- IIPHONORPRIORITY = <u>YES</u>
 - Parm since z/OS 1.8 (2007)



ZAAPAWMT & ZIIPAWMT

- Need for help checking frequency
 - HIPERDISPATCH=NO
 - 12 ms default
 - HIPERDISPATCH=YES
 - 3.2 ms default
 - Automatically increased likelihood to request Help

- Increasing the value(s)
 - Help requested after all zAAPs busy for a longer time
 - Delays GCPs helping
 - Reduces 4HRA impact
 - Reduces overhead of waking up to check for work when low utilization
 - Specialty work waits longer

Workload level GCP overflow controls

Create a Service Class

Service Class Name .	•	•	•	•	•	JZOSLOW
Description	•	•	•	•	•	
Workload Name	•	•	•	•	•	CB (name or ?)
Base Resource Group	•	•	•	•	•	(name or ?)
Cpu Critical	•	•	•	•	•	NO (YES or NO)
I/O Priority Group .	•	•	•	•	•	NORMAL (NORMAL or HIGH
Honor Priority	•					NO (DEFAULT or NO)

Specify BASE GOAL information. Action Codes: I=Insert new period, E=Edit period, D=Delete period. Prevent specialty overflow to GCPs by

- WLM Service
 class
 - Periods not supported

berrea.	Component	PTF
	R7AJ PSY	UA91371
	R7A0 PSY	UA91369
APAR OA50760 –	R79J PSY	UA91372
March 29 th , 2017	R790 PSY	UA91370
March 29 th , 2017	R79J PSY R790 PSY	UA91372 UA91370



Hiperdispatch – zIIP Parking

						сри А	CTII	ΥΤΙΥ	
		z/03 V1R11		SYSTEM I	D 22		START (09/11/200	9-02.30.00
				RPT VERS	ION VIR11	RMF	END (09/11/200	9-03.00.00
CPU	2097	MODEL 737	H/W MODEL	E56 SEQU	ENCE CODE	000000000	00699F1	F HIPE	RDISPATCH=YES
C	PU		TIME	*		LOG PF	20C	I/O I	NTERRUPTS
NUM	TYPE	ONLINE	LPAR BUSY	MVS BUSY	PARKED	SHARE	8	RATE	VIA TPI
0	CP	100.00	96.60	96.74	0.00	100.0	HIGH	1593	2.64
1	CP	100.00	97.51	97.69	0.00	100.0	HIGH	1607	2.73
2	CP	100.00	96.02	96.23	0.00	96.0	MED	5.12	29.30
3	CP	100.00	39.26	80.81	51.23	0.0	LOW	0.00	0.00
4	CP	100.00	48.71	79.90	38.77	0.0	LOW	0.00	0.00
5	CP	100.00	41.06	79.34	48.01	0.0	LOW	0.00	0.00
6	CP	100.00	12.42	78.35	84.11	0.0	LOW	0.00	0.00
7	CP	100.00	0.00		100.00	0.0	LOW	0.00	0.00
8	CP	100.00	0.00		100.00	0.0	LOW	0.00	0.00
9	CP	100.00	33.05	80.34	58.68	0.0	LOW	199.6	1.01
TOTA	L/AVERA	GE	46.46	89.73		296.0		3405	2.62
А	IIP	100.00	57.35	88.68	0.00	32.0	MED		
В	IIP	100.00	46.71	92.85	17.56	0.0	LOW		
С	IIP	100.00	45.27	90.82	17.79	0.0	LOW		
D	IIP	100.00	53.81	85.00	0.00	0.0	LOW		
TOTA	L/AVERA	GE	50.78	89.09		32.0			

- Overflow when <u>online</u> zIIPs are busy
- Unpark > 80%
- Park < 66%
- Focus on VH/VM avoid wasted VLs







How to tune to reduce zIIP usage

- Overflow Reduction
 - Disable those that can wait
 - Make others wait longer
- Tune consumers to reduce overflow
 - JAVA (GC, Application tuning, etc.),
 - DB2 sequential pre-fetch waste



Prevent Overflow ?

Consider

- JAVA batch = wait
- Inefficient zIIP users –
 Disable in product?
- DB2 Test ? = No zIIP

Avoid

- DB2 Production as never use
- Monitors
 - if can't disable in product
 - May cause delay / inaccuracies in sampling



DB2 Sequential Prefetch tuning

- Sequential prefetch Table scans and table space scans
 - Change DB IDX or SQL to reduce # scans
 - Prefetch in DB2 STC so can't control w/ overflow w/ WLM
 - w/o impacting everyone (none zIIP eligible)
 - Reduce prefetch duration & reduce zIIP busy when # < #DB2+1</p>
 - Move table w/ large scan / prefetch to smaller buffer pool
 - Reduces prefetch size and associated zIIP or overflow CPU
 - Reduce % of buffer pool allocated to sequential reads
 - Default 80%
 - VPSEQT = 0 disables pre-fetch



DB2 Utilities

- Run outside of 4HRA or demand peaks
- Reorg fewer databases based on need, not schedule
- Use more efficient DB2 Utilities
- Use reorg while open outside of 4HRA peak





- 125% of 2
 zIIP or 62%
 each
- Overflow to GCP > 50% of a GCP



JAVA – zIIP overflow



- Normalized
 - Using 3 zIIPs
- zIIP Eligible
 - Overflow73% of aGCP
- Non Eligible
 - 61% of aGCP







Garbage Collection Modes

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SMF 121.1 Jobname Max record	GC Data		Retrieval F Start date End Date	ilters	Sta	art Time d Time							≠
Start Time	Uptime Seconds	GC cnt	GC_freed Memory	Used_memor managed_by	GC_Elapsed Time	GC_MODE	PEAK THRDCNT	LIVE THRDCNT	GC_compact Total	System Name	JVM Type	Asid	
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Tuesda	Tuesday 6 th November										
Start	End	Stream	Room	Title	Speaker						
11:45	12:45	IMS	Wellington B	The No Cost Way to Manage the IMS Catalog	David Schipper						
15:00	16:00	IMS	Wellington B	Current Trends in IMS Analytics	David Schipper						
16:30	17:30	zCMPA	Woodcote	zIIP stealing GCP MSUs for Capacity Management	Donald Zeunert						

Wednesday 7 th November									
Start	End	Stream	Room	Title	Speaker				
09:30	10:30	Db2	Nurburgring	Know your onions when it comes to Db2 indexes	Randy Bright				
09:30	10:30	IMS	Wellington B	IMS Checkpoint Pacing	David Schipper				
10:45	11:45	zCMPA	Woodcote	How many GCP MSU is my CF stealing?	Don Zeunert				



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- Paper feedback forms are also available from the Chair person
- Session is LI



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